

Unit 15

Reduction of State Tables / State Assignment

FSM Design Flow

Statement of Problem



State Graph



State Table



Reduction of States



State Assignment



Choice of F/F



Derivation of

F/F Input Equation and

Z Output Equation



Circuit Realization and

Timing Chart



Outline

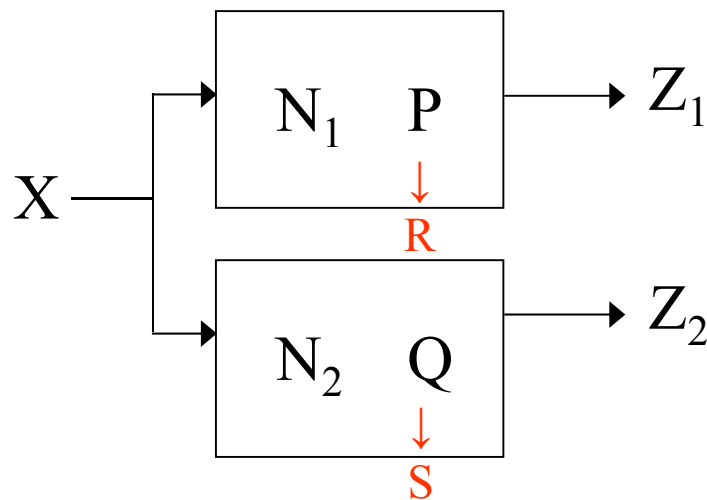
- Equivalent states
- Elimination of redundant states
- Determination of state equivalence using an implication table
- Equivalent sequential circuits
- Incompletely specified state tables
- Derivation of F/F input equations
- Equivalent state assignments
- Guidelines for state assignment
- Using a One-Hot state assignment

Equivalent States

Equivalent states “ \equiv ”

2 machines: N_1, N_2

$\left\{ \begin{array}{l} \text{state P in } N_1 \\ \text{state Q in } N_2 \end{array} \right.$



$P \equiv Q$ iff $R \equiv S$ and $Z_1 = Z_2$

$\lambda(P, x) = \lambda(Q, x)$: output

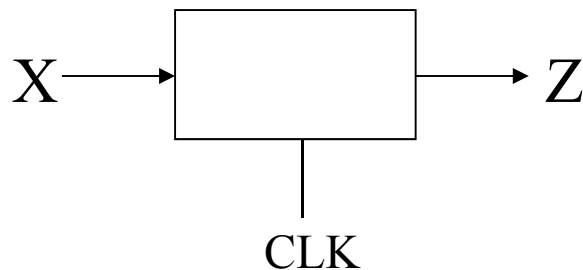
$\delta(P, x) = \delta(Q, x)$: next state

x : any single input

Elimination of Redundant States (1/5)

When setting up states, some extra states may be included
 \Rightarrow eliminate these states

A previous example

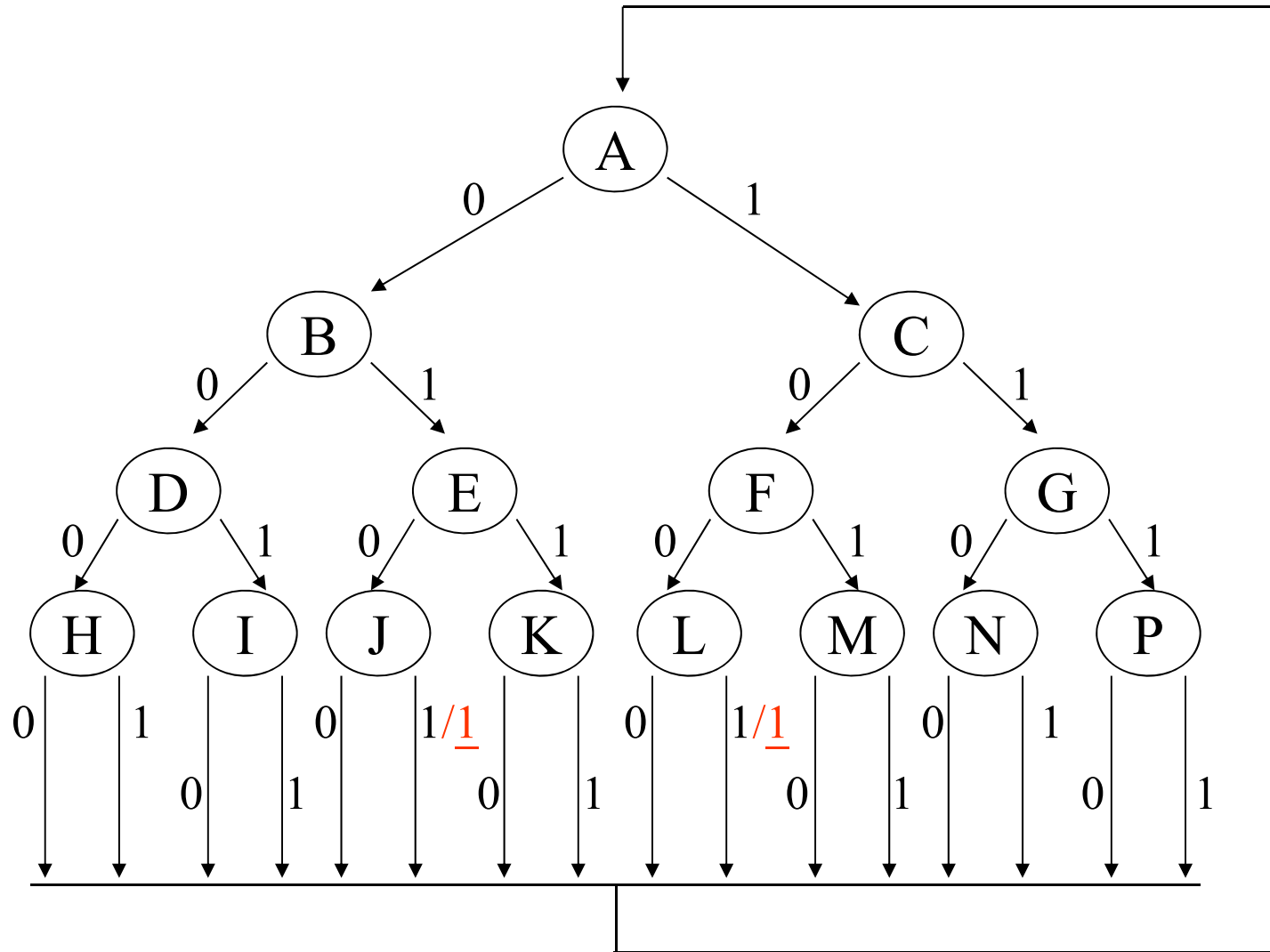


Check 4 consecutive inputs
as a group , then reset
 $Z = 1$ when $X = 0101$ or 1001

A brute force approach:

Reset state A, checks 3 consequent bits of every possible combinations.
After the 4th bit coming-in, give output and reset to state A

Elimination of Redundant States (2/5)





Elimination of Redundant States (3/5)

Input Sequence	P.S.	N.S.		Present Output	
		x = 0	x = 1	x = 0	x = 1
Reset	A	B	C	0	0
0	B	D	E	0	0
1	C	F	G	0	0
00	D	H	I	0	0
01	E	J	K	0	0
10	F	L	M	0	0
11	G	N	P	0	0
000	H	A	A	0	0
001	I	A	A	0	0
010	J	A	A	0	1
011	K	A	A	0	0
100	L	A	A	0	1
101	M	A	A	0	0
110	N	A	A	0	0
111	P	A	A	0	0

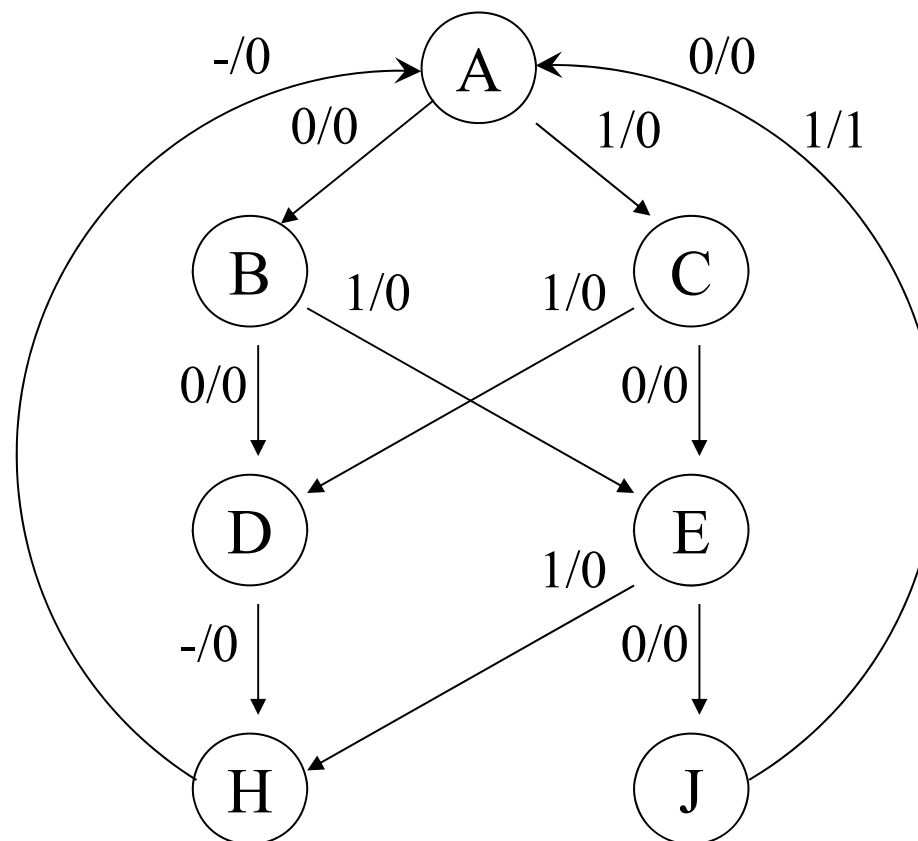
Elimination of Redundant States (4/5)

Equivalent states: the same N.S. & output

Input Sequence	P.S.	N.S.		Present Output	
		x = 0	x = 1	x = 0	x = 1
Reset	A	B	C	0	0
0	B	D	E	0	0
1	C	E F	D G	0	0
00	D	H	H I	0	0
01	E	J	H K	0	0
→ 10	E F	J L	H M	0	0
→ 11	D G	H N	H P	0	0
000	H	A	A	0	0
→ 001	H I	A	A	0	0
010	J	A	A	0	1
→ 011	H K	A	A	0	0
→ 100	J L	A	A	0	1
→ 101	H M	A	A	0	0
→ 110	H N	A	A	0	0
→ 111	H P	A	A	0	0

Elimination of Redundant States (5/5)

P.S.	N.S.		Output	
	x = 0	x = 1	x = 0	x = 1
A	B	C	0	0
B	D	E	0	0
C	E	D	0	0
D	H	H	0	0
E	J	H	0	0
H	A	A	0	0
J	A	A	0	1



Determination of State Equivalence by Using an Implication Table (1/4)



Ex:

P.S.	N.S.		Z
	x = 0	x = 1	
a	d a	c	0
b	f	h	0
c	e c	d a	1
d	a	e	0
e	c	a	1
f	f	b	1
g	b	h	0
h	c	g	1

Determination of State Equivalence by Using an Implication Table (2/4)



Implication chart construction

1. State comparison by implied pair (only for the same-output pair)

b	d - f c - h						
c	X	X					
d	a - d c - e	a - f e - h	X				
e	X	X	a - d c - e	X			
f	X	X	b - d e - f	X	a - b c - f		
g	b - d c - h	b - f	X	a - b e - h	X	X	
h	X	X	d - g c - e	X	a - g	b - g c - f	X
	a	b	c	d	e	f	g

Output shall be the same

N.S.

P.S.	x = 0	x = 1	Z
a	d	c	0
b	f	h	0
c	e	d	1
d	a	e	0
e	c	a	1
f	f	b	1
g	b	h	0
h	c	g	1

Determination of State Equivalence by Using an Implication Table (3/4)



2. Check implied pair iteratively

$a \equiv d$
 $c \equiv e$

b	d-f c-h					
c	X	X				
d	a-d c-e	a-f e-h	X			
e	X	X	a-d c-e	X		
f	X	X	b-d e-f	X	a-b c-f	
g	b-d c-h	b-f	X	a-b e-h	X	X
h	X	X	d-g c-e	X	a-g b-f c-f	X
	a	b	c	d	e	f

P.S.	N.S.		Z
	x = 0	x = 1	
a	d a	c	0
b	f	h	0
c	e c	d a	1
d	a	e	0
e	c	a	1
f	f	b	1
g	b	h	0
h	c	g	1

Determination of State Equivalence by Using an Implication Table (4/4)



N.S.

P.S.	X = 0	X = 1	Z
a	a	c	0
b	f	h	0
c	c	a	1
f	f	b	1
g	b	h	0
h	c	g	1

$a \equiv d$
 $c \equiv e$

Equivalent Sequential Circuits (1/3)

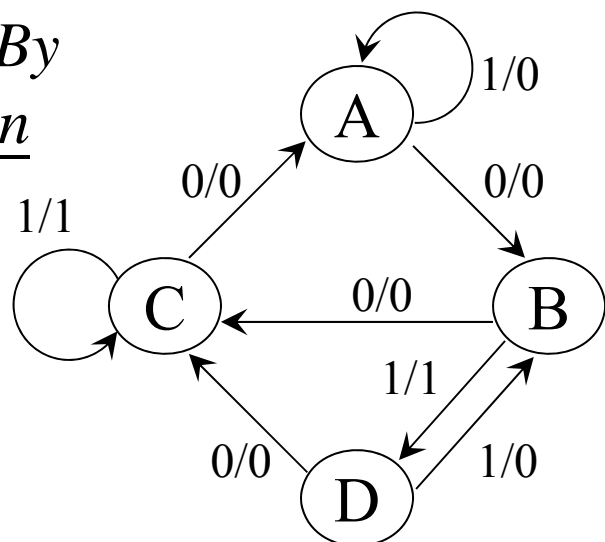
Definition : $N_1 \equiv N_2$ if P in $N_1 \equiv Q$ in N_2
 or S in $N_2 \equiv T$ in N_1

Ex:

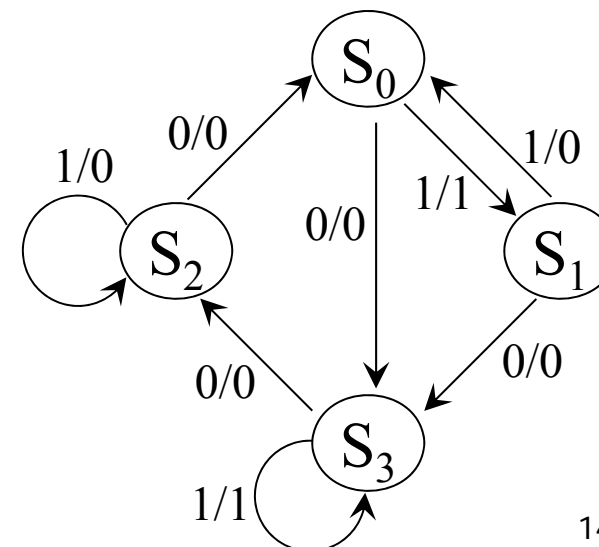
	N_1			
	x = 0	x = 1	x = 0	x = 1
A	B	A	0	0
B	C	D	0	1
C	A	C	0	1
D	C	B	0	0

	N_2			
	x = 0	x = 1	x = 0	x = 1
S_0	S_3	S_1	0	1
S_1	S_3	S_0	0	0
S_2	S_0	S_2	0	0
S_3	S_2	S_3	0	1

Method 1: By observation



$A \equiv S_2$
 $B \equiv S_0$
 $C \equiv S_3$
 $D \equiv S_1$



Equivalent Sequential Circuits (2/3)

		N_1			
		x = 0	x = 1	x = 0	x = 1
A	B	A	0	0	
B	C	D	0	1	
C	A	C	0	1	
D	C	B	0	0	

		N_2			
		x = 0	x = 1	x = 0	x = 1
S_0	S_3	S_1	0	1	
S_1	S_3	S_0	0	0	
S_2	S_0	S_2	0	0	
S_3	S_2	S_3	0	1	

Method 2: By implication table

1. Construct implication table by listing all state pairs. X out pairs with different outputs

S_0	 	C - S_3 D - S_1	A - S_3 C - S_1	
S_1	B - S_3 A - S_0	 	 	C - S_3 B - S_0
S_2	B - S_0 A - S_2	 	 	C - S_0 B - S_2
S_3	 	C - S_2 D - S_3	A - S_2 C - S_3	
	A	B	C	D

Equivalent Sequential Circuits (3/3)

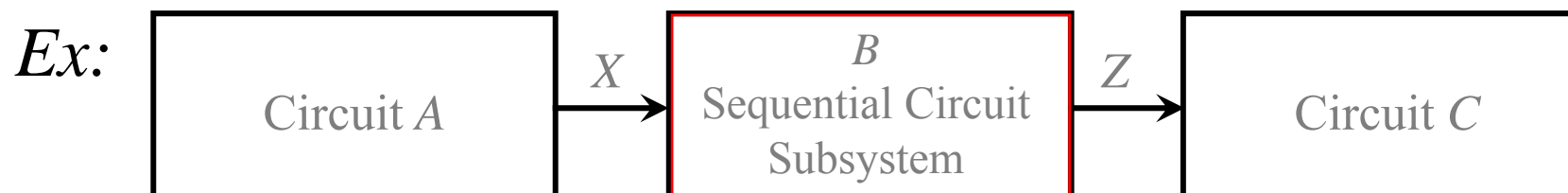
S_0		C - S_3 D - S_1	A - S_3 C - S_1	
S_1	B - S_3 A - S_0			C - S_3 B - S_0
S_2	B - S_0 A - S_2			C - S_0 B - S_2
S_3		C - S_2 D - S_3	A - S_2 C - S_3	
	A	B	C	D

2. *X out additional nonequivalent state pairs*

S_0		C - S_3 D - S_1	A - S_3 C - S_1	
S_1	B - S_3 A - S_0			C - S_3 B - S_0
S_2	B - S_0 A - S_2			C - S_0 B - S_2
S_3		C - S_2 D - S_3	A - S_2 C - S_3	
	A	B	C	D

$$\Rightarrow A \equiv S_2 \quad B \equiv S_0 \quad C \equiv S_3 \quad D \equiv S_1$$

Incompletely Specified State Tables (1/3)



A can only give 100, 110

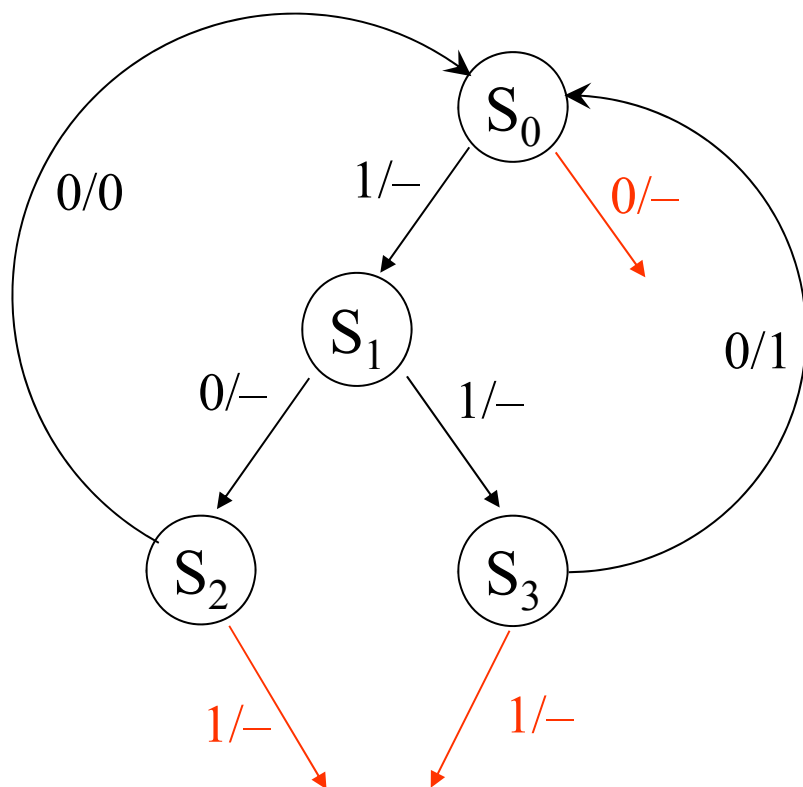
B gives $\left\{ \begin{array}{l} \text{"0"} \\ \text{"1"} \end{array} \right.$ if $\left\{ \begin{array}{l} 100 \text{ received} \\ 110 \text{ received} \end{array} \right.$ at the 3rd bit

	t_0	t_1	t_2		t_0	t_1	t_2
$X =$	1	0	0	$Z =$	—	—	0
	1	1	0		—	—	1

If don't cares are present, the state table is called *incompletely specified*

Incompletely Specified State Tables (2/3)

State Graph

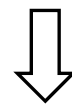


State Table

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
S ₀	—	S ₁	—	—
S ₁	S ₂	S ₃	—	—
S ₂	S ₀	—	0	—
S ₃	S ₀	—	1	—

Incompletely Specified State Tables (3/3)

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
S ₀	– (S ₀)	S ₁	– (0)	–
S ₁	S₂ S ₀	S ₃	– (1)	–
→ S ₂	S ₀	– (S ₁)	0	–
→ S ₃	S ₀	– (S ₃)	1	–



States reduction

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
S ₀	S ₀	S ₁	0	–
S ₁	S ₀	S ₁	1	–



Derivation of F/F Input Equations (1/22)

State Graph

⇒ State Table

⇒ State Assignment

⇒ State Transition Table

Choice of F/F!

⇒ K-Map

⇒ F/F Input Equations



Derivation of F/F Input Equations (2/22)

Type of F/F	Input	Q=0		Q=1	
		Q+=0	Q+=1	Q+=0	Q+=1
D F/F	D	0	1	0	1

Derivation of F/F Input Equations (3/22)

Next State Map

		Q	
	AB	0	1
00		0	1
01		1	0
11		0	0
10		1	×

Q^+

D F/F Input Map

		Q	
	AB	0	1
00		0	1
01		1	0
11		0	0
10		1	×

$$D = Q'A'B' + QB' + AB'$$



Derivation of F/F Input Equations (4/22)

Type of F/F	Input	Q=0		Q=1	
		Q+=0	Q+=1	Q+=0	Q+=1
D F/F	D	0	1	0	1
T F/F	T	0	1	1	0

Derivation of F/F Input Equations (5/22)

Next State Map

		Q	
	AB	0	1
00		0	1
01		1	0
11		0	0
10		1	×

Q^+

T F/F Input Map

		Q	
	AB	0	1
00		0	0
01		1	1
11		0	1
10		1	×

$$T = A'B + AB' + QB$$



Derivation of F/F Input Equations (6/22)

Type of F/F	Input	Q=0		Q=1	
		Q+=0	Q+=1	Q+=0	Q+=1
D F/F	D	0	1	0	1
T F/F	T	0	1	1	0
S-R F/F	S	0	1	0	x
	R	x	0	1	0

Derivation of F/F Input Equations (7/22)

Next State Map

		Q	
	AB	0	1
00		0	1
01		1	0
11		0	0
10		1	×

$$Q^+$$

S-R F/F Input Map

		Q	
	AB	0	1
00		0	×
01		1	0
11		0	0
10		1	×

$$S = AB' + Q'A'B$$

		Q	
	AB	0	1
00		×	0
01		0	1
11		×	1
10		0	×

$$R = QB$$



Derivation of F/F Input Equations (8/22)

Type of F/F	Input	Q=0		Q=1	
		Q+ = 0	Q+ = 1	Q+ = 0	Q+ = 1
D F/F	D	0	1	0	1
T F/F	T	0	1	1	0
S-R F/F	S	0	1	0	x
	R	x	0	1	0
J-K F/F	J	0	1	x	x
	K	x	x	1	0

Derivation of F/F Input Equations (9/22)

Next State Map

		Q	
	AB	0	1
00		0	1
01		1	0
11		0	0
10		1	×

Q^+

J-K F/F Input Map

		Q	
	AB	0	1
00		0	×
01		1	×
11		0	×
10		1	×

$J = A'B + AB'$

		Q	
	AB	0	1
00		×	0
01		×	1
11		×	1
10		×	×

$K = B$

Derivation of F/F Input Equations (10/22)

Ex :

P.S.	N.S.		Z	
	x=0	x=1	x=0	x=1
S_0	S_1	S_2	0	0
S_1	S_3	S_2	0	0
S_2	S_1	S_4	0	0
S_3	S_5	S_2	0	0
S_4	S_1	S_6	0	0
S_5	S_5	S_2	1	0
S_6	S_1	S_6	0	1

7 states \Rightarrow 3 F/Fs

\Rightarrow A, B, C

Let $S_0 = 000$, $S_1 = 110$,
 $S_2 = 001$, $S_3 = 111$,
 $S_4 = 011$, $S_5 = 101$,
 $S_6 = 010$



Derivation of F/F Input Equations (11/22)

Transition table

A	B	C	A ⁺ B ⁺ C ⁺		Z	
			x = 0	x = 1	x = 0	x = 1
0	0	0	110	001	0	0
1	1	0	111	001	0	0
0	0	1	110	011	0	0
1	1	1	101	001	0	0
0	1	1	110	010	0	0
1	0	1	101	001	1	0
0	1	0	110	010	0	1



Derivation of F/F Input Equations (12/22)

(1) Using D F/F

ABC	D_A		D_B		D_C	
	0	1	0	1	0	1
000	1	0	1	0	0	1
110	1	0	1	0	1	1
001	1	0	1	1	0	1
111	1	0	0	0	1	1
011	1	0	1	1	0	0
101	1	0	0	0	1	1
010	1	0	1	1	0	0

ABC	D_A		D_B		D_C	
	0	1	0	1	0	1
000	1	0	1	0	0	1
001	1	0	1	1	0	1
010	1	0	1	1	0	0
011	1	0	1	1	0	0
100	×	×	×	×	×	×
101	1	0	0	0	1	1
110	1	0	1	0	1	1
111	1	0	0	0	1	1

Derivation of F/F Input Equations (13/22)

ABC	D_A		D_B		D_C	
	$X=0$	$X=1$	$X=0$	$X=1$	$X=0$	$X=1$
000	1	0	1	0	0	1
001	1	0	1	1	0	1
010	1	0	1	1	0	0
011	1	0	1	1	0	0
100	×	×	×	×	×	×
101	1	0	0	0	1	1
110	1	0	1	0	1	1
111	1	0	0	0	1	1

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	1	0	0
11	1	1	0	0
10	1	1	0	0

$$A^+ = D_A = X'$$

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	0	0	1
11	1	0	0	1
10	1	1	0	1

$$B^+ = D_B = X'C' + A'C + A'B$$

BC \ XA	00	01	11	10
00	0	×	×	1
01	0	1	1	1
11	0	1	1	0
10	0	1	1	0

$$C^+ = D_C = A + XB'$$

Q^+	D
0	0
1	1



Derivation of F/F Input Equations (14/22)

(2) Choose S-R F/F

ABC	$A^+B^+C^+$ X=		$S_A R_A$ X=		$S_B R_B$ X=		$S_C R_C$ X=		Q	Q ⁺	S	R
	0	1	0	1	0	1	0	1				
000	1	0	0	1	0	1	0	1	0	0	0	×
110	1	1	×	0	0	1	:	:	0	1	1	0
001	1	1	1	0	0	×	:	:	1	0	0	1
111	1	0	×	0	0	1	:	:	1	1	×	0
011	1	1	1	0	0	×	:	:				
101	1	0	×	0	0	1	:	:				
010	1	1	1	0	0	×	:	:				



Derivation of F/F Input Equations (15/22)

(3) Using J-K F/F

ABC	$A^+B^+C^+$		$J_A K_A$		$J_B K_B$		$J_C K_C$		Q	Q ⁺	J	K
	0	1	0	1	0	1	0	1	0	0	0	×
000	110	001	1×	0×	1×	0×	0×	1×	0	0	0	×
110	111	001	×	0	×	1	×	0	×	1	×	×
001	110	011	1×	1×	:	:	:	:	1	0	×	1
111	101	001	×	0	×	1	:	:	1	1	×	0
011	110	010	1×	×	0	:	:	:	:	:	:	:
101	101	001	×	0	0×	:	:	:	:	:	:	:
010	110	010	×	0	×	0	:	:	:	:	:	:

Derivation of F/F Input Equations (16/22)

Next State Maps

ABC	A+		B+		C+	
	0 X=	1	0 X=	1	0 X=	1
000	1	0	1	0	0	1
001	1	0	1	1	0	1
010	1	0	1	1	0	0
011	1	0	1	1	0	0
100	×	×	×	×	×	×
101	1	0	0	0	1	1
110	1	0	1	0	1	1
111	1	0	0	0	1	1

BC	XA			
	00	01	11	10
00	1	×	×	0
01	1	1	0	0
11	1	1	0	0
10	1	1	0	0

BC	XA			
	00	01	11	10
00	1	×	×	0
01	1	0	0	1
11	1	0	0	1
10	1	1	0	1

BC	XA			
	00	01	11	10
00	0	×	×	1
01	0	1	1	1
11	0	1	1	0
10	0	1	1	0

Derivation of F/F Input Equations (17/22)

Q	Q ⁺	J	K
0	0	0	×
0	1	1	×
1	0	×	1
1	1	×	0

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	1	0	0
11	1	1	0	0
10	1	1	0	0

$$A^+$$

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	×	×	0
11	1	×	×	0
10	1	×	×	0

$$J_A = X'$$

BC \ XA	00	01	11	10
00	×	×	×	×
01	×	0	1	×
11	×	0	1	×
10	×	0	1	×

$$K_A = X$$

Derivation of F/F Input Equations (18/22)

Q	Q ⁺	J	K
0	0	0	×
0	1	1	×
1	0	×	1
1	1	×	0

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	0	0	1
11	1	0	0	1
10	1	1	0	1

$$B^+$$

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	0	0	1
11	×	×	×	×
10	×	×	×	×

$$J_B = X'A' + A'C$$

BC \ XA	00	01	11	10
00	×	×	×	×
01	×	×	×	×
11	0	1	1	0
10	0	0	1	0

$$K_B = AC + XA$$

Derivation of F/F Input Equations (19/22)

Q	Q ⁺	J	K
0	0	0	×
0	1	1	×
1	0	×	1
1	1	×	0

BC \ XA	00	01	11	10
00	0	×	×	1
01	0	1	1	1
11	0	1	1	0
10	0	1	1	0

$$C^+$$

BC \ XA	00	01	11	10
00	0	×	×	1
01	×	×	×	×
11	×	×	×	×
10	0	1	1	0

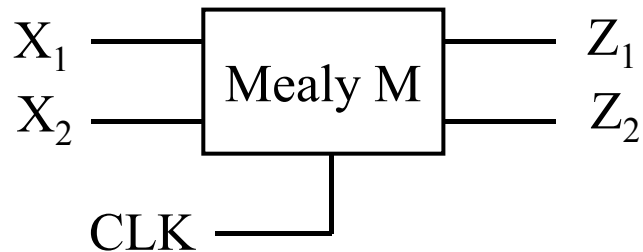
$$J_C = A + XB'$$

BC \ XA	00	01	11	10
00	×	×	×	×
01	1	0	0	0
11	1	0	0	1
10	×	×	×	×

$$K_C = X'A' + A'B$$

Derivation of F/F Input Equations (20/22)

Ex: Another example with 2 inputs X_1, X_2 , 2 outputs Z_1, Z_2



P.S.	N.S.				Z			
	$x_1x_2=00$	$x_1x_2=01$	$x_1x_2=11$	$x_1x_2=10$	$x_1x_2=00$	$x_1x_2=01$	$x_1x_2=11$	$x_1x_2=10$
S_0	S_0	S_0	S_1	S_1	00	00	01	01
S_1	S_1	S_3	S_2	S_1	00	10	10	00
S_2	S_3	S_3	S_2	S_2	11	11	00	00
S_3	S_0	S_3	S_2	S_0	00	00	00	00

4 states \Rightarrow 2 F/Fs A, B

State Assignment: $S_0 = 00, S_1 = 01, S_2 = 11, S_3 = 10$

Derivation of F/F Input Equations (21/22)

Transition table

P.S. AB	N.S.				Z			
	x_1x_2	x_1x_2	x_1x_2	x_1x_2	x_1x_2	x_1x_2	x_1x_2	x_1x_2
00	00	00	01	01	00	00	01	01
01	01	10	11	01	00	10	10	00
11	10	10	11	11	11	11	00	00
10	00	10	11	00	00	00	00	00

Derivation of F/F Input Equations (22/22)

D F/F

		X_1X_2			
	AB	00	01	11	10
00	0	0	0	0	0
01	0	1	1	0	0
11	1	1	1	1	0
10	0	1	1	0	0

$$D_A = A^+ = X_2B + AB + X_2A$$

		X_1X_2			
	AB	00	01	11	10
00	0	0	1	1	0
01	1	0	1	1	0
11	0	0	1	1	0
10	0	0	1	0	0

$$D_B = B^+ = X_1A' + X_2'A'B + X_1B + X_1X_2$$

		X_1X_2			
	AB	00	01	11	10
00	0	0	0	0	0
01	0	1	1	0	0
11	1	1	0	0	0
10	0	0	0	0	0

$$Z_1 = X_2A'B + X_1'AB$$

		X_1X_2			
	AB	00	01	11	10
00	0	0	1	1	0
01	0	0	0	0	0
11	1	1	0	0	0
10	0	0	0	0	0

$$Z_2 = X_1A'B' + X_1'AB$$

Equivalent State Assignments (1/3)

State Graph

⇒ State Table

⇒ State Assignment

⇒ State Transition Table

← (Choice of F/F)

⇒ K-Map

⇒ F/F Input Equations

Different state assignments

⇒ different F/F Input Equations !

Equivalent State Assignments (2/3)

Cost of logic strongly depend on state assignments

Ex:

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
	AB	AB		
S ₁	00	11	0	0
S ₂	01	10	0	1
S ₃	10	01	1	0

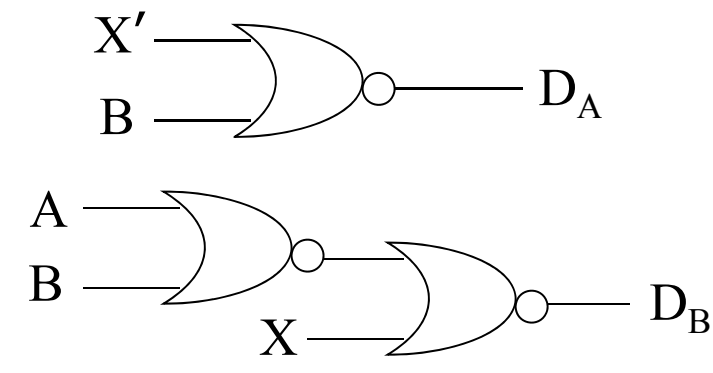
* Assignment 1

S₁ = 00 , S₂ = 01 , S₃ = 10

D F/F:2 F/Fs A, B

$$D_A = XB' \quad D_B = X'(A+B)$$

NOR implementation



Equivalent State Assignments (3/3)

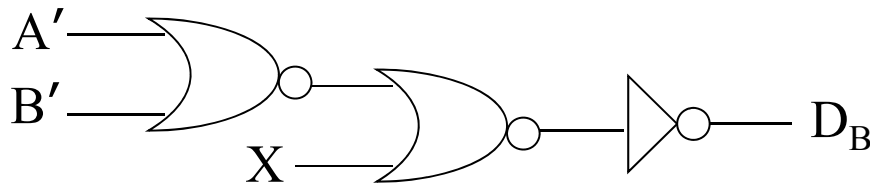
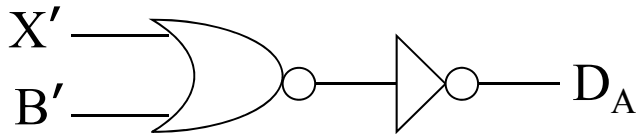
* Assignment 2

$$S_1 = 11, S_2 = 10, S_3 = 01$$

D F/F:2 F/Fs A, B

$$D_A = X' + B' \quad D_B = X + A'B'$$

NOR implementation



\Rightarrow 2 more gates for
NOR implementation



Guidelines for State Assignment (1/16)

Adjacent states: differ in only one variable

2 states:

$S_1(0\ 1\ 0)$	$S_2(0\ 1\ 1)$	} adjacent (S_1, S_2)
$(1\ 0\ 0)$	$(1\ 1\ 0)$	
$S_1(0\ 1\ 0)$	$S_2(1\ 1\ 1)$	} not adjacent
$(0\ 0\ 1)$	$(1\ 1\ 1)$	

Guidelines for State Assignment (2/16)



1. States which have the same next state should be given adjacent assignment
2. States which are next states of the same state should be given adjacent assignment
3. States which have the same output should be given adjacent assignment

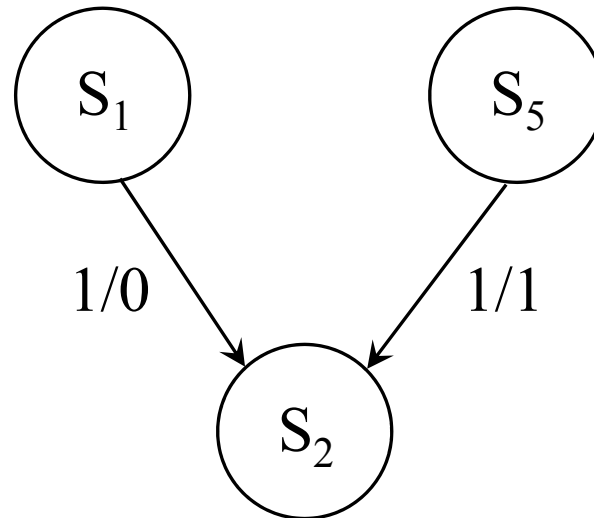
To minimize
output function

In above order !!

No guarantee a minimum solution

Guidelines for State Assignment (3/16)

Guideline 1 : states have the same n.s. with the same input

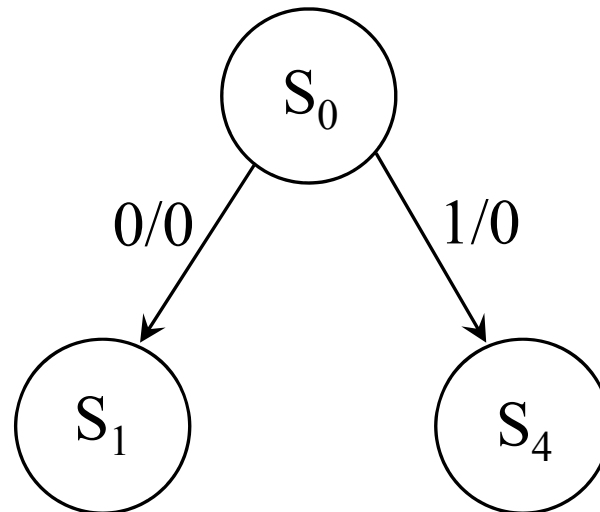


S_1, S_5 are adjacent

Guidelines for State Assignment (4/16)



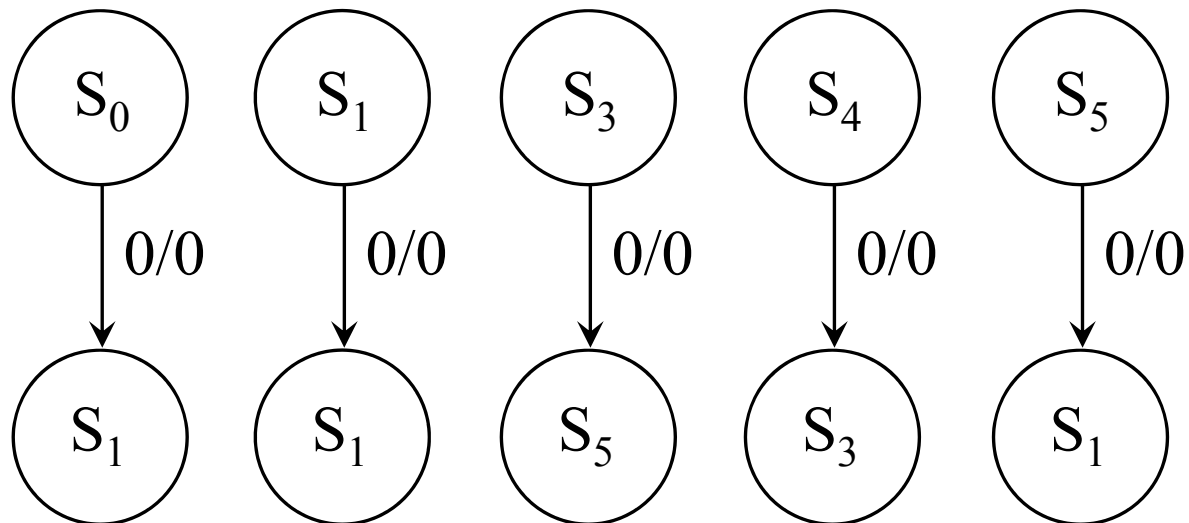
Guideline 2 : states have the same p.s.



S_1, S_4 are adjacent

Guidelines for State Assignment (5/16)

Guideline 3 : states have the same output with the same input



S_0, S_1, S_3, S_4, S_5 are adjacent

Guidelines for State Assignment (6/16)

1. List relationship

Ex:

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
S_0	S_1	S_2	0	0
S_1	S_3	S_2	0	0
S_2	S_1	S_4	0	0
S_3	S_5	S_2	0	0
S_4	S_1	S_6	0	0
S_5	S_5	S_2	1	0
S_6	S_1	S_6	0	1

3 F/Fs A, B, C

States having the same n.s.

Rule 1 : $(S_0, S_2, S_4, S_6)(S_3, S_5)(S_0, S_1, S_3, S_5)(S_4, S_6)$

Rule 2 : $(S_1, S_2)(S_2, S_3)(S_1, S_4)(S_2, S_5) \times 2 (S_1, S_6) \times 2$

States having the same p.s.

Guidelines for State Assignment (7/16)

2. Use an assignment map

	A		
		0	1
BC			
00		S ₀	
01		S ₂	S ₅
11		S ₄	S ₃
10		S ₆	S ₁

Assignment #1

	A		
		0	1
BC			
00		S ₀	
01		S ₁	S ₆
11		S ₃	S ₄
10		S ₅	S ₂

Assignment #2

States having the same n.s.

Rule 1 : (S₀, S₂, S₄, S₆)(S₃, S₅)(S₀, S₁, S₃, S₅)(S₄, S₆)

Rule 2 : (S₁, S₂)(S₂, S₃)(S₁, S₄)(S₂, S₅) × 2 (S₁, S₆) × 2

States having the same p.s.

High
priority



Guidelines for State Assignment (8/16)

Assignment #1	P.S.	N.S.		Z	
		x = 0	x = 1	x = 0	x = 1
000	S ₀	S ₁	S ₂	0	0
110	S ₁	S ₃	S ₂	0	0
001	S ₂	S ₁	S ₄	0	0
111	S ₃	S ₅	S ₂	0	0
011	S ₄	S ₁	S ₆	0	0
101	S ₅	S ₅	S ₂	1	0
010	S ₆	S ₁	S ₆	0	1

6 gates, 13 literals



Guidelines for State Assignment (9/16)

Another straight forward assignment (binary)

	P.S.	N.S.		Z	
		x = 0	x = 1	x = 0	x = 1
000	S ₀	S ₁	S ₂	0	0
001	S ₁	S ₃	S ₂	0	0
010	S ₂	S ₁	S ₄	0	0
011	S ₃	S ₅	S ₂	0	0
100	S ₄	S ₁	S ₆	0	0
101	S ₅	S ₅	S ₂	1	0
110	S ₆	S ₁	S ₆	0	1

10 gates, 39 literals



Guidelines for State Assignment (10/16)

Why guideline help ?

ABC	A+B+C+					Z	
	x=0		x=1			x=0	x=1
000 S ₀	110 S ₁	001 S ₂	001 S ₂	001 S ₂	0	0	
110 S ₁	111 S ₃	001 S ₂	001 S ₂	001 S ₂	0	0	
001 S ₂	110 S ₁	011 S ₄	011 S ₄	011 S ₄	0	0	
111 S ₃	101 S ₅	001 S ₂	001 S ₂	001 S ₂	0	0	
011 S ₄	110 S ₁	010 S ₆	010 S ₆	010 S ₆	0	0	
101 S ₅	101 S ₅	001 S ₀	001 S ₀	001 S ₀	1	0	
010 S ₆	110 S ₁	010 S ₆	010 S ₆	010 S ₆	0	1	

State transition table

		A+B+C+				Z	
		x=0		x=1		x=0	x=1
S ₀	000	110 S ₁	001 S ₂	001 S ₂	001 S ₂	0	0
S ₂	001	110 S ₁	011 S ₄	011 S ₄	011 S ₄	0	0
S ₆	010	110 S ₁	010 S ₆	010 S ₆	010 S ₆	0	1
S ₄	011	110 S ₁	010 S ₆	010 S ₆	010 S ₆	0	0
–	100	×××	×××	×××	×××	–	–
S ₅	101	101 S ₅	001 S ₂	001 S ₂	001 S ₂	1	0
S ₁	110	111 S ₃	001 S ₂	001 S ₂	001 S ₂	0	0
S ₃	111	101 S ₅	001 S ₂	001 S ₂	001 S ₂	0	0

Reordering →



Guidelines for State Assignment (11/16)

BC \ XA	00	01	11	11
00	S_1 S_0	×	×	S_2 S_0
01	S_1 S_2	S_5 S_5	S_2 S_5	S_4 S_2
11	S_1 S_4	S_5 S_3	S_2 S_3	S_6 S_4
10	S_1 S_6	S_3 S_1	S_2 S_1	S_6 S_6

$S_0 S_2 S_4 S_6$

$S_3 S_5$

$S_1 S_3 S_5$

$S_4 S_6$

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	1	0	0
11	1	1	0	0
10	1	1	0	0

$A^+ = X'$

BC \ XA	00	01	11	10
00	1	×	×	0
01	1	0	0	1
11	1	0	0	1
10	1	1	0	1

BC \ XA	00	01	11	10
00	0	×	×	1
01	0	1	1	1
11	0	1	1	0
10	0	1	1	0

due to $S_2 S_5$ assignment



Guidelines for State Assignment (12/16)

Ex:

P.S.	N.S.		Z	
	x = 0	x = 1	x = 0	x = 1
a	a	c	0	0
b	d	f	0	1
c	c	a	0	0
d	d	b	0	1
e	b	f	1	0
f	c	e	1	0

Guidelines

1. $(b,d)(c,f)(b,e)$

2. $(a,c) \times 2, (d,f)(b,d)$

$(b,f)(c,e)$

3. $(a,c)(b,d)(e,f)$

Guidelines for State Assignment (13/16)

Assignment #1

	Q_1		
		0	1
$Q_2 Q_3$			
00	a	c	$a = 000$
01		e	$b = 011$
11	b	d	$c = 100$
10		f	$d = 111$
			$e = 101$
			$f = 110$

Assignment #2

	Q_1		
		0	1
$Q_2 Q_3$			
00	c	a	$a = 100$
01		e	$b = 111$
11	d	b	$c = 000$
10	f		$d = 011$
			$e = 101$
			$f = 010$

1. $(\underline{b,d})(\underline{c,f})(\underline{b,e})$
2. $(\underline{a,c}) \times 2, (\underline{d,f})(\underline{b,d})$
 $(\underline{b,f})(\underline{c,e})$
3. $(\underline{a,c})(\underline{b,d})(\underline{e,f})$

1. $(\underline{b,d})(\underline{c,f})(\underline{b,e})$
2. $(\underline{a,c}) \times 2, (\underline{d,f})(\underline{b,d})$
 $(\underline{b,f})(\underline{c,e})$
3. $(\underline{a,c})(\underline{b,d})(\underline{e,f})$



Guidelines for State Assignment (14/16)

For assignment #2

	P.S.	N.S.		Z	
	$Q_1 Q_2 Q_3$	$x = 0$	$x = 1$	$x = 0$	$x = 1$
m_4	100	100	000	0	0
m_7	111	011	010	0	1
m_0	000	000	100	0	0
m_3	011	011	111	0	1
m_5	101	111	010	1	0
m_2	010	000	101	1	0
m_1	001	---	---	-	-
m_6	110	---	---	-	-

Guidelines for State Assignment (15/16)

State transition table

	P.S. $Q_1Q_2Q_3$	N.S.		Z	
		$x=0$	$x=1$	$x=0$	$x=1$
m_0	000	000	100	0	0
m_1	001	---	---	-	-
m_2	010	000	101	1	0
m_3	011	011	111	0	1
m_4	100	100	000	0	0
m_5	101	111	010	1	0
m_6	110	---	---	-	-
m_7	111	011	010	0	1

next state map

$Q_2Q_3 \backslash XQ_1$	00	01	11	10
00	000	100	000	100
01	×	111	010	×
11	011	011	010	111
10	000	×	×	101

Guidelines for State Assignment (16/16)

D F/F Implementation

		XQ ₁			
		00	01	11	10
Q ₂ Q ₃	00		1		1
	01	×	1		×
	11				1
	10		×	×	1

$$Q_1^+ = X'Q_1Q_2' + XQ_1'$$

		XQ ₁			
		00	01	11	10
Q ₂ Q ₃	00				
	01	×	1	1	×
	11	1	1	1	1
	10		×	×	

$$Q_2^+ = Q_3$$

		XQ ₁			
		00	01	11	10
Q ₂ Q ₃	00				
	01	×	1		×
	11	1	1		1
	10		×	×	1

$$Q_3^+ = X'Q_3 + XQ_1'Q_2$$

		XQ ₁			
		00	01	11	10
Q ₂ Q ₃	00				
	01	×	1		×
	11			1	1
	10	1	×	×	

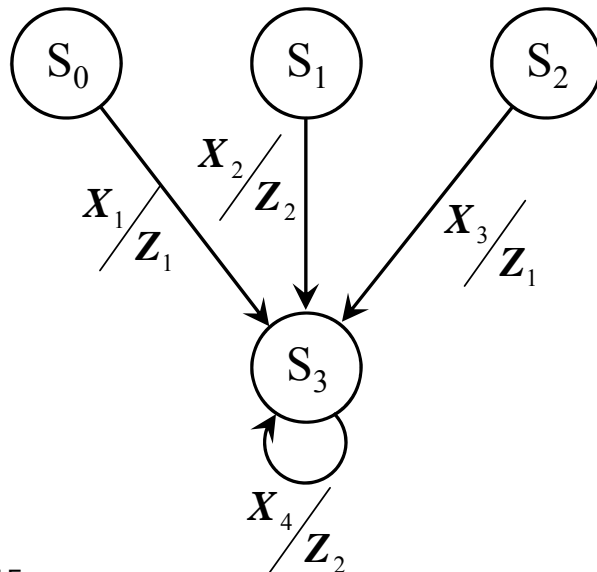
$$Z = X'Q_2'Q_3 + X'Q_2Q_3' + XQ_2Q_3$$

Using a One-Hot State Assignment (1/2)

- Design with FPGAs, not to minimize the number of F/F
- Reduce the number of logic cells and interconnections between cells to minimize delay

One-hot state assignment

- Using one F/F only for each state
- 4 states use 4 F/Fs (Q_0, Q_1, Q_2, Q_3) with following state assignment
 $S_0: Q_0Q_1Q_2Q_3=1000, S_1: 0100, S_2: 0010, S_3: 0001$



Write next-state and output eq.
directly by inspecting the state graph

$$Q_3^+ = X_1Q_0 + X_2Q_1 + X_3Q_2 + X_4Q_3$$

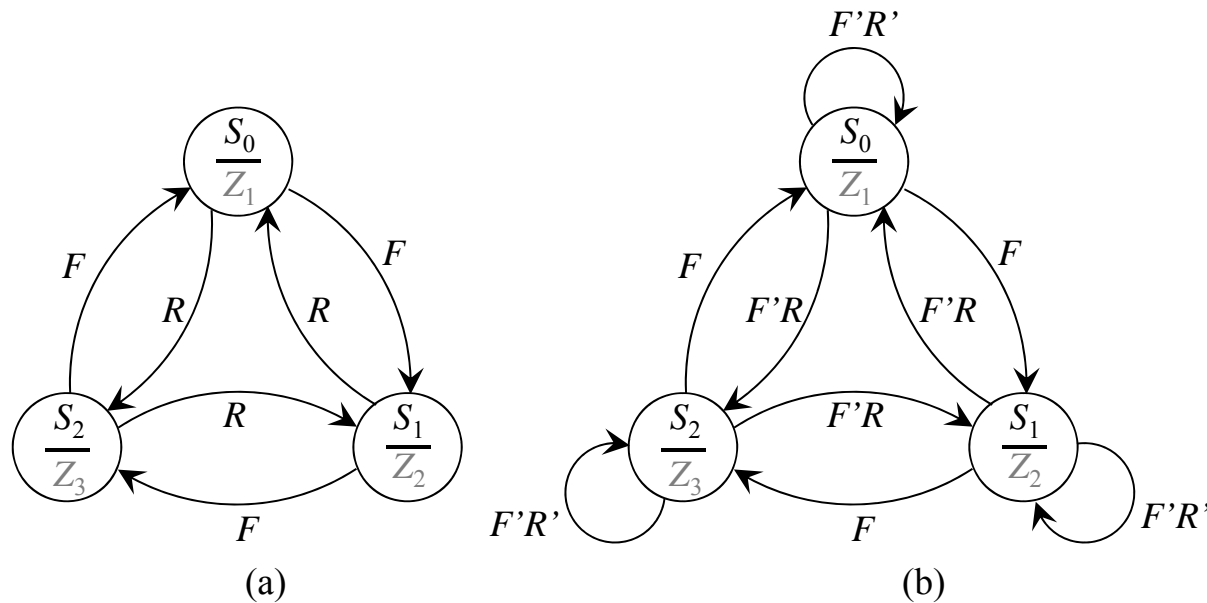
$$Z_1 = X_1Q_0 + X_3Q_2$$

$$Z_2 = X_2Q_1 + X_4Q_3$$

Using a One-Hot State Assignment (2/2)

- For Moore machine,

$$S_0: Q_0Q_1Q_2=100, S_1: 010, S_2: 001$$



$$Q_0^+ = F'R'Q_0 + FQ_2 + F'RQ_1$$

$$Q_1^+ = F'R'Q_1 + FQ_0 + F'RQ_2$$

$$Q_2^+ = F'R'Q_2 + FQ_1 + F'RQ_0$$

$$Z_1 = Q_0 \quad Z_2 = Q_1 \quad Z_3 = Q_2$$